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# Understanding and Comparing Unsuccessful Executions in Large Datacenters

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#### Abstract

The project aims at comparing two different traces coming from large datacenters, focusing in particular on unsuccessful executions of jobs and tasks submitted by users. The objective of this project is to compare the resource waste caused by unsuccessful executions, their impact on application performance, and their root causes. We show the strong negative impact on CPU and RAM usage and on task slowdown. We analyze patterns of unsuccessful jobs and tasks, particularly focusing on their interdependency. Moreover, we uncover their root causes by inspecting key workload and system attributes such asmachine locality and concurrency level.

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## 1 Introduction

In today's world there is an ever growing demand for efficient, large scale computations. The rising trend of "big data" put the need for efficient management of large scaled parallelized computing at an all time high. This fact also increases the demand for research in the field of distributed systems, in particular in how to schedule computations effectively, avoid wasting resources and avoid failures.

In 2011 Google released a month long data trace of their own cluster management system[1] *Borg*, containing a lot of data regarding scheduling, priority management, and failures of a real production workload. This data was 2009 This data was the foundation of the 2015 Rosá et al. paper *Understanding the Dark Side of Big Data Clusters: An Analysis beyond Failures*[2], which in its many conclusions highlighted the need for better cluster management highlighting the high amount of failures found in the traces.

In 2019 Google released an updated version of the *Borg* cluster traces[3], not only containing data from a far bigger workload due to improvements in computational technology, but also providing data from 8 different *Borg* cells from datacenters located all over the world. These new traces are therefore about 100 times larger than the old traces, weighing in terms of storage spaces approximately 8TiB (when compressed and stored in JSONL format)[4], requiring a considerable amount of computational power to analyze them and the implementation of special data engineering techniques for analysis of the data.

This project aims to repeat the analysis performed in 2015 to highlight similarities and differences in workload this decade brought, and expanding the old analysis to understand even better the causes of failures and how to prevent them. Additionally, this report provides an overview of the data engineering techniques used to perform the queries and analyses on the 2019 traces.

Cluster	Timezone
А	America/New York
В	America/Chicago
С	America/New York
D	America/New York
Е	Europe/Helsinki
F	America/Chicago
G	Asia/Singapore
Н	Europe/Brussels

Figure 1. Approximate geographical location obtained from the datacenter's timezone of each cluster in the 2019 Google Borg traces.

## 2 State of the art

In 2015, Dr. Andrea Rosà et al. published a research paper titled *Understanding the Dark Side of Big Data Clusters: An Analysis beyond Failures*[2] in which they performed several analysis on unsuccessful executions in the Google's 2011 Borg cluster traces with the aim of identifying their resource waste, their impacts on the performance of the application, and any causes that may lie behind such failures. The salient conclusion of that research is that actually lots of computations performed by Google would eventually end in failure, then leading to large amounts of computational power being wasted.

However, with the release of the new 2019 traces, the results and conclusions found by that paper could be potentially outdated in the current large-scale computing world. The new traces not only provide updated data on Borg's workload, but provide more data as well: the new traces contain data from 8 different Borg "cells" (i.e. clusters) in datacenters across the world, from now on referred as "Cluster A" to "Cluster H".

The geographical location of each cluster can be consulted in Figure 1. The information in that table was provided by the 2019 traces documentation[4].

The new 2019 traces provide richer data even on a cluster by cluster basis. For example, the amount and variety of server configurations per cluster increased significantly from 2011. An overview of the machine configurations in the cluster analyzed with the 2011 traces and in the 8 clusters composing the 2019 traces can be found in Figure 2. Additionally, in Figure 3, the same machine configuration data is provided for the 2019 traces providing a cluster-by-cluster distribution of the machines.

CPU	RAM	% Machines	CPU	RAM	% Machines	CPU	RAM	% Machines
0.5	0.5	53.47%	0.25	0.25	0.99%	0.5	0.97	0.03%
0.5	0.25	30.74%	0.5	0.12	0.43%	1	0.5	0.02%
0.5	0.75	7.95%	0.5	0.03	0.04%	0.5	0.06	0.01%
1	1	6.32%						

#### (a) 2011 data

CPU (NCU)	RAM (NMU)	Machine count	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines
Unknown	Unknown	8729	1.639218%	0.386719	0.166748	27011	5.072393%	0.958984	0.250000	600	0.112674%
1.000000	0.500000	124234	23.329891%	1.000000	1.000000	12286	2.307187%	0.500000	0.062500	54	0.010141%
0.591797	0.333496	103013	19.344801%	0.591797	0.166748	9902	1.859496%	0.500000	0.250000	34	0.006385%
0.259277	0.166748	78078	14.662260%	1.000000	0.250000	7550	1.417814%	0.479492	0.250000	12	0.002253%
0.708984	0.333496	55801	10.478864%	0.958984	1.000000	3552	0.667030%	0.708984	0.250000	6	0.001127%
0.386719	0.333496	36237	6.804943%	0.259277	0.333496	3024	0.567877%	0.591797	0.250000	4	0.000751%
0.958984	0.500000	31151	5.849843%	0.591797	0.666992	1000	0.187790%	0.708984	0.500000	2	0.000376%
0.708984	0.666992	29594	5.557454%	0.259277	0.083374	634	0.119059%	0.479492	0.500000	2	0.000376%

#### (b) 2019 data

**Figure 2.** Overview of machine configurations in term of CPU and Memory power in 2011 and 2019 (all clusters aggregated) traces. In the 2019 traces NCU stands for "Normalized Compute Unit" and NMU stands for "Normalized Compute Unit": both are [0, 1] normalizations of resource values. While memory was measured in terms of capacity, CPU power was measured in "Google Compute Units" (GCUs), an opaque umbrella metric used by Google that factors in CPU clock, number of cores/processors, and CPU ISA architecture.

CPU (NCU)	RAM (NMU)	Machine count	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines
Unknown	Unknown	1377	1.623170%	Unknown	Unknown	134	0.264812%	Unknown	Unknown	1466	2.274208%
0.591797	0.333496	29487	34.758469%	0.591797	0.333496	16184	31.982926%	0.259277	0.166748	15754	24.439204%
1.000000	0.500000	13440	15.842705%	1.000000	0.500000	9790	19.347061%	0.386719	0.333496	11104	17.225652%
0.708984	0.333496	12495	14.728764%	0.708984	0.333496	8448	16.694992%	0.591797	0.333496	10404	16.139741%
0.386719	0.333496	9057	10.676144%	0.958984	0.500000	5502	10.873088%	0.958984	0.500000	6634	10.291334%
0.386719	0.166748	5265	6.206238%	0.708984	0.666992	3832	7.572823%	1.000000	0.500000	5654	8.771059%
0.708984	0.666992	4608	5.431784%	1.000000	1.000000	2214	4.375321%	0.386719	0.166748	3580	5.553660%
1.000000	1.000000	4446	5.240823%	0.591797	0.166748	2152	4.252796%	0.708984	0.666992	2900	4.498774%
0.591797	0.166748	2484	2.928071%	0.386719	0.333496	816	1.612584%	1.000000	1.000000	2736	4.244361%
0.958984	0.500000	1143	1.347337%	0.958984	1.000000	618	1.221296%	1.000000	0.250000	2132	3.307375%
0.958984	1.000000	654	0.770917%	0.591797	0.666992	500	0.988103%	0.958984	1.000000	766	1.188297%
1.000000	0.250000	366	0.431431%	0.386719	0.166748	412	0.814197%	0.708984	0.333496	620	0.961807%
0.479492	0.250000	6	0.007073%					0.958984	0.250000	600	0.930781%
0.708984	0.250000	6	0.007073%					0.591797	0.166748	112	0.173746%

(a) A cluster

#### (b) Cluster B

(c) Cluster C

CPU (NCU)	RAM (NMU)	Machine count	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines
Unknown	Unknown	498	0.794309%	Unknown	Unknown	536	0.671915%	Unknown	Unknown	1432	2.299958%
0.591797	0.333496	28394	45.288376%	0.259277	0.166748	38452	48.202377%	1.000000	0.500000	41340	66.396839%
0.386719	0.333496	8402	13.401174%	0.708984	0.333496	11786	14.774608%	0.708984	0.333496	6878	11.046866%
0.259277	0.166748	8020	12.791885%	0.958984	0.500000	8646	10.838389%	0.591797	0.333496	5564	8.936430%
0.386719	0.166748	5806	9.260559%	0.708984	0.666992	7606	9.534674%	0.958984	0.500000	2172	3.488484%
0.708984	0.666992	4380	6.986092%	1.000000	0.500000	5586	7.002457%	0.386719	0.166748	1544	2.479843%
0.708984	0.333496	3924	6.258772%	0.386719	0.166748	4470	5.603470%	0.708984	0.666992	1244	1.998008%
0.591797	0.166748	2548	4.064055%	0.259277	0.333496	1268	1.589530%	1.000000	0.250000	792	1.272044%
0.259277	0.333496	426	0.679469%	0.259277	0.083374	634	0.794765%	0.958984	1.000000	536	0.860878%
1.000000	0.500000	292	0.465739%	0.591797	0.333496	324	0.406158%	0.386719	0.333496	398	0.639234%
0.591797	0.250000	4	0.006380%	1.000000	0.250000	268	0.335957%	1.000000	1.000000	344	0.552504%
0.708984	0.500000	2	0.003190%	1.000000	1.000000	138	0.172993%	0.500000	0.250000	18	0.028910%
-				0.500000	0.062500	54	0.067693%				
				0.500000	0.250000	4	0.005014%				

(d) Cluster D

(e) Cluster E

(f) Cluster F

Machine count					
	% Machines	CPU (NCU)	RAM (NMU)	Machine count	% Machines
1566	2.261568%	Unknown	Unknown	1720	2.933251%
15852	22.892958%	1.000000	0.500000	36324	61.946178%
11808	17.052741%	0.591797	0.333496	4826	8.230158%
7968	11.507134%	0.708984	0.333496	3682	6.279205%
7830	11.307839%	0.958984	0.500000	2858	4.873973%
4690	6.773150%	0.386719	0.333496	2596	4.427163%
4258	6.149269%	1.000000	1.000000	2030	3.461919%
4196	6.059731%	1.000000	0.250000	1892	3.226577%
3864	5.580267%	0.386719	0.166748	1244	2.121491%
2606	3.763503%	0.708984	0.666992	766	1.306320%
2100	3.032754%	0.591797	0.666992	500	0.852689%
1330	1.920744%	0.958984	1.000000	200	0.341076%
778	1.123563%				
378	0.545896%				
12	0.017330%				
6	0.008665%				
2	0.002888%				
	12 6 2 Cluster G	12 0.017330% 6 0.008665% 2 0.002888% Cluster G	12         0.017330%           6         0.008665%           2         0.002888%	12 0.017330% 6 0.008665% 2 0.002888% Cluster G (h) C	12         0.017330%           6         0.008665%           2         0.002888%           (h) Cluster H

Figure 3. Overview of machine configurations in terms of CPU and RAM resources for each cluster in the 2019 traces. Refer to figure 2 for a column legend.

Type code	Description
EVICT	The job or task was terminated in order to free computational resources for an
	higher priority job
FAIL	The job or task terminated its execution unsuccesfully due to a failure
FINISH	The job or task terminated succesfully
KILL	The job or task terminated its execution because of a manual request to stop it

Figure 4. Overview of job and task termination event types.



Figure 5. Typical transitions between task/job event types according to Google

## 3 Background information

*Borg* is Google's own cluster management software able to run thousands of different jobs. Among the various cluster management services it provides, the main ones are: job queuing, scheduling, allocation, and deallocation due to higher priority computations.

The core structure of Borg is a cell, a set of machines usually all within the same cluster, whose work is allocated by the same cluster-management system and hence a cell is handled as a unit. Each cell may run large computational workload that is submitted to Borg. Such workload is called "job", which outlines the computations that a user wants to run and is made up of several "tasks". A task is an executable program, consisting of multiple processes, which has to be run on a single machine. Those tasks may be ran sequentially or in parallel, and the condition for a job's successful termination is nontrivial.

#### 3.1 Traces

The data relative to the events happening while Borg cell processes the workload is then encoded and stored as rows of several tables that make up a single usage trace. Such data comes from the information obtained by the cell's management system and single machines that make up the cell. Each table is identified by a key, usually a timestamp.

In general events can be of two kinds, there are events that are relative to the status of the schedule, and there are other events that are relative to the status of a task itself.

Figure 5 shows the expected transitions between event types.

#### 3.2 Traces contents

The traces provided by Google contain mainly a collection of job and task events spanning a month of execution of the 8 different clusters. In addition to this data, some additional data on the machines' configuration in terms of resources (i.e. amount of CPU and RAM) and additional machine-related metadata.

Due to Google's policy, most identification related data (like job/task IDs, raw resource amounts and other text values) were obfuscated prior to the release of the traces. One obfuscation that is noteworthy in the scope of this thesis is related to CPU and RAM amounts, which are expressed respetively in NCUs (*Normalized Compute Units*) and NMUs (*Normalized Memory Units*).

NCUs and NMUs are defined based on the raw machine resource distributions of the machines within the 8 clusters. A machine having 1 NCU CPU power and 1 NMU memory size has the maximum amount of raw CPU power and raw RAM size found in the clusters. While RAM size is measured in bytes for normalization purposes, CPU power was measured in GCU (*Google Compute Units*), a proprietary CPU power measurement unit used by Google that combines several parameters like number of processors and cores, clock frequency, and architecture (i.e. ISA).

### 3.3 Overview of traces' format

The traces have a collective size of approximately 8TiB and are stored in a Gzip-compressed JSONL (JSON lines) format, which means that each table is represented by a single logical "file" (stored in several file segments) where each carriage return separated line represents a single record for that table.

There are namely 5 different table "files":

machine\_configs, which is a table containing each physical machine's configuration and its evolution over time;

instance\_events, which is a table of task events;

- collection\_events, which is a table of job events;
- machine\_attributes, which is a table containing (obfuscated) metadata about each physical machine and its evolution over time;

**instance\_usage**, which contains resource (CPU/RAM) measures of jobs and tasks running on the single machines.

The scope of this thesis focuses on the tables machine\_configs, instance\_events and collection\_events.

#### 3.4 Remark on traces size

While the 2011 Google Borg traces were relatively small, with a total size in the order of the tens of gigabytes, the 2019 traces are quite challenging to analyze due to their sheer size. As stated before, the traces have a total size of 8 TiB when stored in the format provided by Google. Even when broken down to table "files", unitary sizes still reach the single tebibyte mark (namely for machine\_configs, the largest table in the trace).

Due to this constraints, a careful data engineering based approach was used when reproducing the 2015 DSN paper analysis. Bleeding edge data science technologies like Apache Spark were used to achieve efficient and parallelized computations. This approach is discussed with further detail in the following section.

# 4 Project Requirements and Analysis Methodology

The aim of this project is to repeat the analysis performed in 2015 on the dataset Google has released in 2019 in order to find similarities and differences with the previous analysis, and ultimately find whether computational power is indeed wasted in this new workload as well. The 2019 data comes from 8 Borg cells spanning 8 different datacenters located in different geographical positions, all focused on computational oriented workloads. The data collection time span matches the entire month of May 2019.

Due to the inherent complexity in analyzing traces of this size, non-trivial data engineering tecniques were adopted to performed the required computations. We used the framework Apache Spark to perform efficient and parallel Map-Reduce computations. In this section, we discuss the technical details behind our approach.

## 4.1 Introduction on Apache Spark

Apache Spark is a unified analytics engine for large-scale data processing. In layman's terms, Spark is really useful to parallelize computations in a fast and streamlined way.

In the scope of this thesis, Spark was used essentially as a Map-Reduce framework for computing aggregated results on the various tables. Due to the sharded nature of table "files", Spark is able to spawn a thread per file and run computations using all processors on the server machines used to run the analysis.

Spark is also quite powerful since it provides automated thread pooling services, and it is able to efficiently store and cache intermediate computation on secondary storage without any additional effort required from the data engineer. This feature was especially useful due to the sheer size of the analyzed data, since the computations required to store up to 1TiB of intermediate data on disk.

The chosen programming language for writing analysis scripts was Python. Spark has very powerful native Python bindings in the form of the *PySpark* API, which were used to implement the various queries.

#### 4.2 Query architecture

#### 4.2.1 Overview

In general, each query written to execute the analysis follows a Map-Reduce template. Traces are first read, then parsed, and then filtered by performing selections, projections and computing new derived fields.

After this preparation phase, the trace records are often passed through a groupby() operation, which by choosing one or many record fields sorts all the records into several "bins" containing records with matching values for the selected fields. Then, a map operation is applied to each bin in order to derive some aggregated property value for each grouping.

Finally, a reduce operation is applied to either further aggregate those computed properties or to generate an aggregated data structure for storage purposes.

#### 4.2.2 Parsing table files

As stated before, table "files" are composed of several Gzip-compressed shards of JSONL record data. The specification for the types and constraints of each record is outlined by Google in the form of a protobuffer specification file found in the trace release package[5]. This file was used as the oracle specification and was a critical reference for writing the query code that checks, parses and carefully sanitizes the various JSONL records prior to actual computations.

The JSONL encoding of traces records is often performed with non-trivial rules that required careful attention. One of these involved fields that have a logically-wise "zero" value (i.e. values like "0" or the empty string). For these values the key-value pair in the JSON object is outright omitted. When reading the traces in Apache Spark is therefore necessary to check for this possibility and insert back the omitted record attributes.

#### 4.2.3 The queries

Most queries use only two or three fields in each trace records, while the original table records often are made of a couple of dozen fields. In order to save memory during the query, a projection is often applied to the data by the means of a .map() operation over the entire trace set, performed using Spark's RDD API.

Another operation that is often necessary to perform prior to the Map-Reduce core of each query is a record filtering process, which is often motivated by the presence of incomplete data (i.e. records which contain fields whose values

is unknown). This filtering is performed using the .filter() operation of Spark's RDD API.

The core of each query is often a groupby() followed by a map() operation on the aggregated data. The groupby() groups the set of all records into several subsets of records each having something in common. Then, each of this small clusters is reduced with a map() operation to a single record. The motivation behind this way of computing data is that for the analysis in this thesis it is often necessary to analyze the behaviour w.r.t. time of either task or jobs by looking at their events. These queries are therefore implemented by groupby()-ing records by task or job, and then map()-ing each set of event records sorting them by time and performing the desired computation on the obtained chronological event log.

Sometimes intermediate results are saved in Spark's parquet format in order to compute and save intermediate results beforehand.

#### 4.3 Query script design

In this section we aim to show the general complexity behind the implementations of query scripts by explaining in detail some sampled scripts to better appreciate their behaviour.

#### 4.3.1 The "task slowdown" query script

One example of analysis script with average complexity and a pretty straightforward structure is the pair of scripts task\_slowdown.py and task\_slowdown\_table.py used to compute the "task slowdown" tables (namely the tables in Figure 9).

"Slowdown" is a task-wise measure of wasted execution time for tasks with a FINISH termination type. It is computed as the total execution time of the task divided by the execution time actually needed to complete the task (i.e. the total time of the last execution attempt, successful by definition).

The analysis requires to compute the mean task slowdown for each task priority value, and additionally compute the percentage of tasks with successful terminations per priority. The query therefore needs to compute the execution time of each execution attempt for each task, determine if each task has successful termination or not, and finally combine this data to compute slowdown, mean slowdown and ultimately the final table found in Figure 9.

Figure 6 shows a schematic representation of the query structure.

The query first starts reading the instance\_events table, which contains (among other data) all task event logs containing properties, event types and timestamps. As already explained in the previous section, the logical table file is actually stored as several Gzip-compressed JSONL shards. This is very useful for processing purposes, since Spark is able to parse and load in memory each shard in parallel, i.e. using all processing cores on the server used to run the queries.

After loading the data, a selection and a projection operation are performed in the preparation phase so as to "clean up" the records and fields that are not needed, leaving only useful information to feed in the "group by" phase. In this query, the selection phase removes all records that do not represent task events or that contain an unknown task ID or a null event timestamp. In the 2019 traces it is quite common to find incomplete records, since the log process is unable to capture the sheer amount of events generated by all jobs in a exact and deterministic fashion.

Then, after the preparation stage is complete, the task event records are grouped in several bins, one per task ID. Performing this operation the collection of unsorted task event types is rearranged to form groups of task events all relating to a single task.

These obtained collections of task events are then sorted by timestamp and processed to compute intermediate data relating to execution attempt times and task termination counts. After the task events are sorted, the script iterates over the events in chronological order, storing each execution attempt time and registering all execution termination types by checking the event type field. The task termination is then equal to the last execution termination type, following the definition originally given in the 2015 Rosá et al. DSN paper.

If the task termination is determined to be unsuccessful, the tally counter of task terminations for the matching task property is increased. Otherwise, all the task termination attempt time deltas are returned. Tallies and time deltas are saved in an intermediate time file for fine-grained processing.

Finally, the task\_slowdown\_table.py processes this intermediate results to compute the percentage of successful tasks per execution and computing slowdown values given the previously computed execution attempt time deltas. Finally, the mean of the computed slowdown values is computed resulting in the clear and coincise tables found in Figure 9.



Figure 6. Diagram of the script used for the "task slowdown" query.

## 5 Analysis: Performance Input of Unsuccessful Executions

Our first investigation focuses on replicating the analysis done by the paper of Rosá et al. paper[2] regarding usage of machine time and resources.

In this section we perform several analyses focusing on how machine time and resources are wasted, by means of a temporal vs. spatial resource analysis from the perspective of single tasks as well as jobs. We then compare the results from the 2019 traces to the ones that were obtained before to understand the workload evolution inside Borg between 2011 and 2019.

We discover that the spatial and temporal impact of unsuccessful executions is very significant, more than in the 2011 traces. In particular, resource usage is overall dominated by tasks with a final KILL termination event.

#### 5.1 Temporal Impact: Machine Time Waste

The goal of this analysis is to understand how much time is spent in doing useless computations by exploring how machine time is distributed over task events and submissions.

Before delving into the analysis itself, we define three kinds of events in a task's lifecycle:

submission: when a task is added or re-added to the Borg system queue, waiting to be scheduled;

**scheduling:** when a task is removed from the Borg queue and its actual execution of potentially useful computations starts;

termination: when a task terminates its computations either successfully or unsuccessfully.

By partitioning the set of all terminating tasks by their termination event, the analysis aims to measure the total time spent by tasks in 3 different execution phases:

- **resubmission time:** the total of all time intervals between every task termination event and the immediately succeding task submission event, i.e. the total time spent by tasks waiting to be resubmitted in Borg after a termination;
- **queue time:** the total of all time intervals between every task submission event and the following task scheduling event, i.e. the total time spent by tasks queuing before execution;
- **running time:** the total of all time intervals between every task scheduling event and the following task termination event, i.e. the total time spent by tasks "executing" (i.e. performing potentially useful computations) in the clusters.

In the 2019 traces, an additional "Unknown" measure is counted. This measure collects all the times in which the event transitions between the register events do not allow to safely assume in which execution phase a task may be. Unknown measures are mostly caused by faults and missed event writes in the task event log that was used to generate the traces.

The analysis results are depicted in Figure 7 as a comparison between the 2011 and 2019 traces, aggregating the data from all clusters. Additionally, in Figure 8 cluster-by-cluster breakdown result is provided for the 2019 traces.

The striking difference between 2011 and 2019 data is in the machine time distribution per task termination type. In the 2019 traces, 94.38% of global machine time is spent on tasks that are eventually KILLed. FINISH, EVICT and FAIL tasks respectively register totals of 4.20%, 1.18% and 0.25% machine time.

Considering instead the distribution between execution phase times, the comparison shows very similar behaviour between the two traces, having the "Running" time being dominant (at a total of 16.63% across task terminations in 2019) over the queue and resubmission phases (with respective totals in 2019 of 3.26% and 0.004%).

However, another noteworthy difference between 2011 and 2019 data lies in the new "Unknown" trace dataset present only in the latter traces, registering a total 80.12% of global machine time across al terminations. This data can be interpreted as a strong indication of the "poor quality" of the 2019 traces w.r.t. of accuracy of task event logging.

Considering instead the behaviour of each single cluster in the 2019 traces, no significant difference beween them can be observed. The only notable difference lies between the "Running time"-"Unknown time" ratio in KILLed tasks, which is at its highest in cluster A (at 30.78% by 58.71% of global machine time) and at its lowest in cluster H (at 8.06% by 84.77% of global machine time).

The takeaway from this analysis is that in the 2019 traces a lot of computation time is wasted in the execution of tasks that are eventually KILLed, i.e. unsuccessful.



**Figure 7.** Relative task time spent in each execution phase w.r.t. task termination in 2011 and 2019 (all clusters aggregated) traces. The x-axis shows task termination type, Y axis shows total time % spent. Colors break down the time in execution phases. "Unknown" execution times are 2019 specific and correspond to event time transitions that are not consider "typical" by Google.



Figure 8. Relative task time spent in each execution phase w.r.t. clusters in the 2019 trace. Refer to Figure 7 for axes description.

#### 5.2 Average Slowdown per Task

This analysis aims to measure the average of an ad-hoc defined parameter we call "slowdown". We define it as the ratio between the total response time across all executions of the task and the response time (i.e. queue time and running time) of the last execution of said task. This metric is especially useful to analyze the impact of unsuccesful executions on each task total execution time w.r.t. the intrinsic workload (i.e. computational time) of tasks.

Refer to Figure 9 for a comparison between the 2011 and 2019 mean task slowdown measures broke down by task priority. Additionally, said means are computed on a cluster-by-cluster basis for 2019 data in Figure 10.

In 2015 Rosá et al.[2] measured mean task slowdown per each task priority value, which at the time were numeric values between 0 and 11. However, in 2019 traces, task priorities are given as a numeric value between 0 and 500. Therefore, to allow an easier comparison, mean task slowdown values are computed by task priority tier over the 2019 data. Priority tiers are semantically relevant priority ranges defined in the Tirmazi et al. 2020[3] that introduced the 2019 traces. Equivalent priority tiers are also provided next to the 2011 priority values in the table covering the 2015 analysis.

In the given tables, the % **finished** column corresponds to the percentage of **FINISHed** tasks for that priority or tier. **Mean response [s] (last execution)** instead shows the mean response time of the last task execution of each task in that priority/tier. **Mean response [s] (all executions)** provides a very similar figure, though this column shows the mean response time across all executions. **Mean slowdown** instead provides the mean slowdown value for each task priority/tier.

Comparing the tables in Figure 9 we observe that the maximum mean slowdown measure for 2019 data (i.e. 7.84, for the BEB tier) is almost double of the maximum measure in 2011 data (i.e. 3.39, for priority 3 corresponding to the BEB tier). The "Best effort batch" tier, as the name suggest, is a lower priority tier where failures are more tolerated. Therefore, due to the increased concurrency in the 2019 clusters compared to 2011 and the higher machine time spent for unsuccesful executions (as observed in the previous analysis) and increase slowdown rate for this class is not particularly surprising.

The amount of non-successful task terminations in the 2019 traces is also rather high when compared to 2011 data, as it can evinced by the low percentage of FINISHed tasks across priority tiers.

Another noteworthy difference is in the mean response times for all and last executions: while the mean response is overall shorter in time in the 2019 traces by an order of magnitude, the new traces show an overall significantly higher mean response time than in the 2011 data.

Across 2019 single clusters (as in Figure 10), the data shows a mostly uniform behaviour, other than for some noteworthy mean slowdown spikes. Indeed, cluster A has 82.97 mean slowdown in the "Free" tier, cluster G has 19.06 and 14.57 mean slowdown in the "BEB" and "Production" tier respectively, and Cluster D has 12.04 mean slowdown in its "Free" tier.

#### 5.3 Spatial Impact: Resource Waste

In this analysis we aim to understand how physical resources of machines in the Borg cluster are used to complete tasks. In particular, we compare how CPU and Memory resource allocation and usage are distributed among tasks based on their termination type.

Due to limited computational resources w.r.t. the data analysis process, the resource usage for clusters E to H in the 2019 traces is missing. However, a comparison between 2011 resource usage and the aggregated resource usage of clusters A to D in the 2019 traces can be found in Figure 11. Additionally, a cluster-by-cluster breakdown for the 2019 data can be found in Figure 12.

From these figures it is clear that, compared to the relatively even distribution of used resources in the 2011 traces, the distribution of resources in the 2019 Borg clusters became strikingly uneven, registering a combined 86.29% of CPU resource usage and 84.86% memory usage for KILLed tasks. Instead, all other task termination types have a significantly lower resource usage: EVICTed, FAILed and FINISHed tasks register respectively 8.53%, 3.17% and 2.02% CPU usage and 9.03%, 4.45%, and 1.66% memory usage. This resource distribution can also be found in the data from individual clusters in Figure 12, with always more than 80% of resources devoted to KILLed tasks.

Considering now requested resources instead of used ones, a comparison between 2011 and the aggregation of all A-H clusters of the 2019 traces can be found in Figure 13. Additionally, a cluster-by-cluster breakdown for single 2019 clusters can be found in Figure 14.

Here KILLed jobs dominate even more the distribution of resources, reaching a global 97.21% of CPU allocation and

2011 priority	Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
0	Erroo	53.80%	2845	1767	3.37
1	Fiee	67.44%	3598	2939	2.58
2		90.78%	1835	1782	1.15
3	Deet	95.62%	9683	8294	3.39
4	Best	78.05%	2006	1890	1.69
5	errort	100%	58	58	1
6	Datch	77.99%	567	567	1.02
8		45.48%	1159	1151	1.01
9	Production	23.35%	504	496	1.07

(a) 2011 data

Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
Best effort batch	11.06%	4139	113	7.84
Free	42.85%	1374	8	1.15
Mid	2.71%	18187	157	2.55
Monitoring	2.74%	834226	130	2.05
Production	13.54%	54789	24	6.68

(b) 2019 data, aggregated

Figure 9. Mean task slowdown for each cluster and each priority "tier" for 2011 and 2019 data. % finished is the percentage of tasks with FINISH termination w.r.t. priority, Mean response [s] (last execution) is the mean response time (queue+execution time, in seconds) for the last task execution w.r.t. priority, Mean response [s] (all executions) is the response time (in seconds) of all executions, Mean slowdown is the mean slowdown measure w.r.t. priority. Priorities with no successfully terminated jobs have been omitted.

Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown	Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
Best effort batch	212.62%	71108	14201	5.17	Best effort batch	71.84%	1018454	550288	8.47
Free	0.33%	5769	1203	82.97	Free	45.21%	12047	5588	1.18
Mid	46.22%	8510	9135	1.16	Mid	8.82%	225147	336262	1.11
Monitoring	2.82%	1200998	1054458	2.86	Monitoring	4.12%	2627612	2024679	1.51
Production	27.21%	4546	16845	4.12	Production	30.92%	182604	466329	9.71
		(a) Cluster	A				(b) Cluster	В	

Tier

Free

Mid

Monitoring

Production

Monitoring

Production

Best effort batch

Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
Best effort batch	52.96%	1236666	997117	7.40
Free	73.36%	172214	5553	1.12
Mid	95.4%	579844	248553	2.04
Monitoring	5.88%	2159459	1761833	1.74
Production	3.61%	352603	357993	4.14

(c) Cluster C

Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
Best effort batch	0.47%	280811	205838	8.06
Free	48.15%	33050	40073	1.44
Mid	0.46%	62123	83322	10.31
Monitoring	37.71%	1415296	1263746	2.82
Production	1.96%	231639	414149	8.54

(e) Cluster E

Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
Best effort batch	44.29%	1368306	1563086	6.14
Free Mid	45.86% 31.36%	187447 200116	37069 110201	1.09 7.60

Mean response [s]

(last execution)

1154060

22831

228762

1588844

279565

(d) Cluster D

2079134 297168

(f) Cluster F

Mean response [s]

(all executions)

1135023

225269

913816

349364

1682711 492372

5506

Mean slowdown

12.04

1.15

2 56

2.16

5.51

2.08

5.94

8.42% 3.65%

% finished

50.56%

42.82%

86.34%

2.21%

6.53%

Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown	Tier	% finished	Mean response [s] (last execution)	Mean response [s] (all executions)	Mean slowdown
Best effort batch	104.33%	294959	184724	19.06	Best effort batch	107.03%	947368	527812	7.33
Free	33.85%	64718	15473	1.14	Free	28.79%	310534	290058	1.12
Mid	49.06%	732532	706124	3.86	Mid	2.18%	338883	197440	6.49
Monitoring	4.36%	1991341	1676276	1.72	Monitoring	4.96%	2309296	1808698	1.94
Production	26.75%	115953	399050	14.57	Production	2.7%	298799	470783	5.80

(g) Cluster G

(h) Cluster H

Figure 10. Mean task slowdown for each cluster and each task priority for single clusters in the 2019 traces. Refer to 9 for a legend of the columns



**Figure 11.** Percentages of CPU and RAM resources used by tasks w.r.t. task termination type in 2011 and 2019 traces (total of clusters A to D). The x axis is the type of resource, y-axis is the percentage of resource used and color represents task termination. Numeric values are displayed below the graph as a table.





Cluster B (used resources) 100 80 6040200 % Memory % CPU 4.85%7.35% 6.31% 2.59% 8.42% 1.23% 86.25 83.00%



Figure 12. Percentages of CPU and RAM resources used by tasks w.r.t. task termination type for clusters A to D in 2019 traces. Refer to figure 11 for plot explaination.





Figure 13. Percentages of CPU and RAM resources requested by tasks w.r.t. task termination type in 2011 and 2019 traces. The x axis is the type of resource, y-axis is the percentage of resource used and color represents task termination. Numeric values are displayed below the graph as a table.



















Figure 14. Percentages of CPU and RAM resources requested by tasks w.r.t. task termination type for in 2019 traces. Refer to figure 13 for plot explaination.

a global 96.89% of memory allocation. Even in allocations, the KILL lead is followed by (in order) EVICTed, FAILed and FINISHed jobs, with respective CPU allocation figures of 2.73%, 0.06% and 0.0012% and memory allocation figures of 3.04%, 0.06% and 0.012%.

Behaviour across clusters (as evinced in Figure 14) in terms of requested resources is pretty homogeneous, with the exception of cluster A having a relatively high 2.85% CPU and 3.42% memory resource requests from EVICTed tasks and cluster E having a noteworthy 1.67% CPU and 1.31% memory resource resquests from FINISHed tasks.

With more than 98% of both CPU and memory resources used by (and more than 99.99% of both CPU and memory resources requested by) non-successful tasks, it is clear the spatial resource waste is high in the 2019 traces.

## 6 Analysis: Patterns of Task and Job Events

This section aims to use some of the tecniques used in section IV of the Rosá et al. paper[2] to find patterns and interpendencies between task and job events by gathering event statistics at those events. In particular, Section 6.1 explores how tasks of the success of a task is inter-correlated with its own event patterns, which Section 6.2 explores even further by computing task success probabilities based on the number of task termination events of a specific type. Finally, Section 6.3 aims to find similar correlations, but at the job level.

The results found the the 2019 traces seldomly show the same patterns in terms of task events and job/task distributions, in particular highlighting again the overall non-trivial impact of KILL events, no matter the task and job termination type.

## 6.1 Unsuccessful Task Event Patterns

In this analysis we compute the distribution of termination events by type at the task-level events, namely EVICT, FAIL, FINISH and KILL termination events.

A comparison of the termination event distribution between the 2011 and 2019 traces is shown in Figure 15. Additionally, a cluster-by-cluster breakdown of the same data for the 2019 traces is shown in Figure 16.

Each table from these figure shows the mean and the 95-th percentile of the number of termination events per task, broke down by task termination. In addition, the table shows the mean number of EVICT, FAIL, FINISH, and KILL for each task event termination.

The first observation we make is that the mean number of events per EVICTed and FAILed tasks increased more than 5-fold (namely from 2.372 to 78.710 and from 3.130 to 24.962 respectively). Also observing the 95-th percentile we can say that the number of events per task has generally increased overall.

As observed in 2011, 2019 Borg tasks have all a multitude of events with different types, with FINISHed tasks experiencing almost always FINISH events and unsuccessful tasks and the same observation holding for KILLed tasks and their KILL events. Differently from the 2011 data, EVICTed tasks seem to experience an high number of KILL events as well (25.795 on average per task, over 78.710 overall events on average). A similar phenomena can be observed with KILLed jobs and their EVICT events (1.876 on average per task with a 8.763 event overall average).

Considering cluster-by-cluster behaviour in the 2019 traces (as reported in Figure 16) the general observations still hold for each cluster, albeit with event count averages having different magnitudes. Notably, cluster E registers the highest per-event average, with FAILed tasks experiencing 111.471 FAIL events out of 112.384.

## 6.2 Conditional Probability of Task Success

In this analysis we measure the conditional probability of task success given a number of specific unsuccessful (i.e. EVICT, FAIL and KILL) events. This analysis was conducted to better understand how a given number of unsuccessful events could affect the termination of the task it belongs to.

Conditional probabilities of each unsuccessful event type are shown in the form of a plot in Figure 17, comparing the 2011 traces with the overall data from the 2019 ones, and in Figure 18, as a cluster-by-cluster breakdown of the same data for the 2019 traces.

In Figure 17 the 2011 and 2019 plots differ in their x-axis: for 2011 data conditional probabilities are computed for a maximum event count t of 30, while for 2019 data are computed for up to 50 events of a specific kind. Nevertheless, another quite striking difference between the two plots can be seen: while 2011 data has relatively smooth decreasing curves for all event types, the curves in the 2019 data almost immediately plateau with no significant change easily observed after 5 events of any kind.

The presence of even one KILL event almost surely causes the corresponding task to terminate in an unsuccessful way: a task with no KILL events has 97.16% probability of success, but tasks with 1 to 5 KILL events have 0.02%, 0.20%, 0.44%, 0.04%, and 0.07% probabilities of success respectively. The same effect can be observed, albeit in a less drastic fashion, for the EVICT and FAIL curves. The EVICT curve has for tasks with 0 to 5 kill events 19.70%, 15.94%, 1.94%, 1.67%, 0.35% and 0.00% success probabilities repectively. The FAIL probability curve has instead 18.55%, 1.79%, 14.49%, 2.08%, 2.40%, and 1.29% success probabilities for the same range.

Considering cluster-to-cluster behaviour in the 2019 traces (as shown in Figure 18), some clusters show quite similar behaviour to the aggregated plot (namely clusters A, F, and H), while some other clusters show very oscillating probability distribution function curves for EVICT and FINISH curves. KILL behaviour is instead homogeneous even

Task	Mean number of events							
terminatio	on Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL			
EVICT	2.372 (5)	2.094	0.259	0.004	0.015			
FAIL	3.130 (8)	0.350	2.700	0.020	0.060			
FINISH	2.516 (4)	0.302	1.175	1.023	0.016			
KILL	1.094 (1)	0.061	0.008	0.011	1.014			
<b>(a)</b> 2011 data								
Task	М	ean numb	er of ever	nts				
termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL			
EVICT	78.710 (342)	52.242	0.673	0.000	25.795			
FAIL	24.962 (26)	0.290	23.635	0.348	0.691			
FINISH	2.962 (2)	0.022	0.012	2.915	0.013			
KILL	8.763 (16)	1.876	0.143	0.003	6.741			

(b)	2019	data
$(\nu)$	201/	uata

**Figure 15.** Mean number of termination events and their distributions per task type between 2011 and 2019 (all clusters aggregated) traces. The tables show an overall mean accompanied by the 95-th percentile of all termination events, followed by the mean of events per event type of each termination event.

Task	N	lean numl	ber of ever	nts		Task	Ν	lean num	ber of ever	nts	
termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL	termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL
EVICT	103.228 (719)	73.694	0.769	0.000	28.766	EVICT	83.018 (394)	64.817	0.240	0.000	17.962
FAIL	11.819 (26)	0.288	11.062	0.002	0.468	FAIL	20.851 (62)	0.518	19.657	0.001	0.675
FINISH	2.185 (1)	0.019	0.004	2.153	0.008	FINISH	2.995 (4)	0.020	0.021	2.943	0.012
KILL	5.963 (11)	2.350	0.214	0.003	3.396	KILL	9.173 (12)	3.351	0.276	0.004	5.541
(a) Cluster A							(b	) Cluster B			
Task	Ν	lean numl	ber of ever	nts		Task	Ν	lean num	ber of even	nts	
termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL	termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL
EVICT	98.437 (444)	73.716	1.813	0.000	22.908	EVICT	76.759 (366)	62.001	0.700	0.000	14.058
FAIL	52.010 (30)	0.773	48.446	2.035	0.756	FAIL	62.314 (62)	0.496	58.968	0.810	2.040
FINISH	2.507 (2)	0.018	0.013	2.471	0.006	FINISH	3.877 (2)	0.059	0.019	3.789	0.010
KILL	5.452 (6)	1.533	0.116	0.004	3.799	KILL	6.795 (6)	1.960	0.151	0.002	4.682
	(c)	) Cluster C					(d	) Cluster D			
Task	Ν	lean numl	per of even	its		Task	Ν	lean num	ber of ever	nts	
termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL	termination	Overall (95 <sup>th</sup> p.)	EVICT	FAIL	FINISH	KILL
EVICT	17.678 (72)	11.781	0.106	0.000	5.791	EVICT	70.146 (114)	23.974	0.192	0.000	45.980
FAIL	112.384 (28)	0.458	111.471	0.000	0.456	FAIL	41.087 (54)	0.279	39.257	0.000	1.550
FINISH	2.029 (2)	0.014	0.008	1.999	0.008	FINISH	3.129 (4)	0.019	0.004	3.008	0.098
KILL	13.505 (64)	1.288	0.057	0.000	12.160	KILL	10.288 (38)	0.384	0.098	0.001	9.804
(e) Cluster E						(f	) Cluster F				
Task	N	lean num	ber of ever	nts		Task	Me	ean numbe	er of even	ts	
Task termination	M Overall (95 <sup>th</sup> p.)	lean num EVICT	ber of even FAIL	nts FINISH	KILL	Task termination	Me Overall (95 <sup>th</sup> p.)	ean numbo EVICT	er of even FAIL	ts FINISH	KILL
Task termination EVICT	M Overall (95 <sup>th</sup> p.) 136.032 (490)	Iean num EVICT 77.429	ber of even FAIL 0.303	nts FINISH 0.000	KILL 58.299	Task termination EVICT	Ma Overall (95 <sup>th</sup> p.) 14.734 (40)	ean numbe EVICT 6.733	er of even FAIL 0.837	ts FINISH 0.000	KILL 7.165
Task termination EVICT FAIL	M Overall (95 <sup>th</sup> p.) 136.032 (490) 8.948 (8)	Iean num EVICT 77.429 0.016	ber of even FAIL 0.303 8.593	nts FINISH 0.000 0.000	KILL 58.299 0.339	Task termination EVICT FAIL	Me Overall (95 <sup>th</sup> p.) 14.734 (40) 41.067 (120)	ean numbo EVICT 6.733 0.600	er of even FAIL 0.837 37.600	ts FINISH 0.000 0.000	KILL 7.165 2.867
Task termination EVICT FAIL FINISH	M Overall (95 <sup>th</sup> p.) 136.032 (490) 8.948 (8) 14.176 (2)	Iean num EVICT 77.429 0.016 0.015	ber of even FAIL 0.303 8.593 0.002	nts FINISH 0.000 0.000 14.154	KILL 58.299 0.339 0.005	Task termination EVICT FAIL FINISH	Me Overall (95 <sup>th</sup> p.) 14.734 (40) 41.067 (120) 3.681 (2)	ean numbe EVICT 6.733 0.600 0.024	er of even FAIL 0.837 37.600 0.014	ts FINISH 0.000 0.000 3.633	KILL 7.165 2.867 0.011
Task termination EVICT FAIL FINISH KILL	M Overall (95 <sup>th</sup> p.) 136.032 (490) 8.948 (8) 14.176 (2) 32.320 (164)	Iean num EVICT 77.429 0.016 0.015 6.909	ber of even FAIL 0.303 8.593 0.002 0.135	nts FINISH 0.000 0.000 14.154 0.000	KILL 58.299 0.339 0.005 25.276	Task termination EVICT FAIL FINISH KILL	Ma Overall (95 <sup>th</sup> p.) 14.734 (40) 41.067 (120) 3.681 (2) 17.976 (98)	ean numbe EVICT 6.733 0.600 0.024 0.633	er of even FAIL 0.837 37.600 0.014 0.170	ts FINISH 0.000 0.000 3.633 0.000	KILL 7.165 2.867 0.011 17.173

**Figure 16.** Mean number of termination events and their distributions per task type for each cluster in the 2019 traces. The tables show an overall mean accompanied by the 95-th percentile of all termination events, followed by a mean of events per event type of each termination event.



**Figure 17.** Conditional probability of task success given a number of specific unsuccessful events observed, i.e., EVICT, FAIL and KILL for 2011 and 2019 (all clusters aggregated) traces. For 2011 data the probability was computed for a maximum event count of 30, while in 2019 it was computed for up to 50 events of a specific kind.



Figure 18. Conditional probability of task success given a number of specific unsuccessful events observed, i.e., EVICT, FAIL and KILL for each cluster in the 2019 traces.

on a single cluster basis.

## 6.3 Unsuccessful Job Event Patterns

The analysis uses very similar techniques to the ones used in Section 6.1, but focusing at the job level instead. The aim is to better understand the task-job level relationship and to understand how task-level termination events can influence the termination state of a job.

A comparison of the analyzed parameters between the 2011 and 2019 traces is shown in Figure 19. Additionally, a cluster-by-cluster breakdown of the same data for the 2019 traces is shown in Figure 20.

Considering the distribution of number of tasks in a job, the 2019 traces show a decrease for the mean figure (e.g. for FAILed jobs, with a mean 60.5 tasks per job in 2011 and a mean 43.126 tasks per job in 2019) and a fluctuation of the 95-th percentile figure (e.g. for FAILed jobs it rose from 110 to 200, but for KILLed job the figure decreased from 400 to 178).

Considering the distribution of the number of task-wise termination events instead, the 2019 traces show values generally one or two orders of magnitude below the ones in 2011. While the behaviour of EVICTed jobs stays the same, FAILed and KILLed jobs show a dramatic difference in the event distribution, with KILL becoming the most popular event task-wise with mean 12.833 and 11.337 task events per job respectively. Finally, the FINISHed job category has a new event distribution too, with FINISH task events being the most popular at 1.778 events per job in the 2019 traces.

The cluster-by-cluster comparison in Figure 20 shows that the number of tasks per job are generally distributed similarly to the aggregated data, with only cluster H having remarkably low mean and 95-th percentiles overall. Event-wise, for EVICTed, FINISHed, and KILLed jobs again the distributions are similar to the aggregated one. For some clusters (namely B, C, and D), the mean number of FAIL and KILL task events for FINISHed jobs is almost the same. Additionally, it is noteworthy that cluster A has no EVICTed jobs.

Job	# of tasks mean.	Mean number of events								
termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL					
EVICT	1 (1)	1	0	0	0					
FAIL	60.5 (110)	139.0	788.5	49.2	9.5					
FINISH	2.7 (1)	0.4	0.1	$5 \cdot 10^{-4}$	2.7					
KILL	86.8 (400)	13.3	20.9	26.9	62.7					
(a) 2011 data										
Job	# of tasks mean.	M	ean num	ber of ever	nts					
termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL					
EVICT	1.000 (1)	1.000	0.000	0.000	0.000					
FAIL	43.126 (200)	0.114	2.300	0.981	12.833					
FINISH	3.074 (2)	0.005	0.153	1.778	0.014					
KILL	53.919 (178)	0.235	0.103	0.288	11.337					

#### (b) 2019 data

**Figure 19.** Mean number of tasks and event distribution per job type for between 2011 and 2019 (all clusters aggregated) traces. The tables show and mean and 95-th percentile for the number of tasks in a job, and additionally show the mean of job-wise total of task termination events.

Job	# of tasks mean.	Μ	ean num	ber of eve	nts	Job	# of tasks mean.	Ν	lean num	ber of eve	nts
termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL	termination	n (95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL
EVICT	-	-	_	-	_	EVICT	1.000 (1)	1.000	0.000	0.000	0.000
FAIL	90.793 (499)	0.695	0.684	0.086	1.850	FAIL	74.368 (374)	2.003	1.994	0.267	4.944
FINISH	1.187 (1)	0.005	0.001	1.073	0.024	FINISH	6.304 (10)	0.022	0.008	2.349	0.013
KILL	16.533 (10)	1.045	0.074	0.461	1.189	KILL	69.853 (234)	1.696	0.158	0.614	3.009
(a) Cluster A						(b)	Cluster B				
Job	# of tasks mean.	Μ	ean num	ber of eve	nts	Job	# of tasks mean.	Ν	lean num	ber of eve	nts
termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL	termination	n (95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL
EVICT	1.000 (1)	1.001	0.000	0.000	0.000	EVICT	1.000 (1)	1.000	0.000	0.000	0.000
FAIL	41.982 (200)	3.484	0.998	0.376	3.998	FAIL	43.356 (250)	6.112	0.949	0.531	6.498
FINISH	1.991 (1)	0.022	0.017	1.565	0.017	FINISH	2.109 (2)	0.268	0.013	1.723	0.019
KILL	110.681 (652)	0.627	0.059	0.656	2.267	KILL	89.648 (283)	1.013	0.054	0.283	3.256
(c) Cluster C						(d)	Cluster D				
Job	# of tasks mean.	М	ean num	ber of eve	nts	Job	# of tasks mean.	N	lean num	ber of eve	nts
termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL	termination	n (95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL
EVICT	1.000 (1)	1.000	0.000	0.000	0.000	EVICT	1.000 (1)	1.000	0.000	0.000	0.000
FAIL	23.081 (25)	0.247	0.666	0.717	1.588	FAIL	17.161 (8)	0.621	0.546	0.426	7.559
FINISH	7.776 (2)	0.019	0.029	1.934	0.021	FINISH	2.941 (2)	0.015	0.051	1.670	0.162
KILL	88.790 (309)	0.706	0.029	0.461	7.572	KILL	103.889 (361)	0.183	0.064	0.417	5.824
(e) Cluster E						(f)	Cluster F				
Job	# of tasks mean.	M	ean numl	ber of ever	nts	Job	# of tasks mean.	Me	an numb	er of even	ts
termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL	termination	(95 <sup>th</sup> p)	EVICT	FAIL	FINISH	KILL
EVICT	1.000 (1)	1.000	0.000	0.000	0.000	EVICT	1.000 (1)	1.000	0.000	0.000	0.000
FAIL	51.835 (250)	0.556	3.335	0.608	20.352	FAIL	20.504 (1)	0.114	2.300	0.981	12.833
FINISH	8.519 (36)	0.002	0.630	1.760	0.005	FINISH	4.278 (14)	0.005	0.153	1.778	0.014
KILL	37.055 (100)	5.687	0.065	0.080	19.166	KILL	11.023 (3)	0.235	0.103	0.288	11.337
	(g) (	Cluster G					(h) (	luster H			

**Figure 20.** Mean number of tasks and event distribution per job type for each cluster in the 2019 traces. The tables show and mean and 95-th percentile for the number of tasks in a job, and additionally show the mean of job-wise total of task termination events.

7 Conclusions, Future Work and Possible Developments

TBD

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